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**A note on hypergraph connectivity
augmentation**

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Abstract

We prove an abstract version of an edge-splitting theorem for directed hypergraphs that appeared in [1], and use this result to obtain min-max theorems on hypergraph augmentation problems that are linked to orientations. These problems include (k, l) -edge-connectivity augmentation of directed hypergraphs, and (k, l) -partition-connectivity augmentation of undirected hypergraphs by uniform hyperedges.

1 Introduction

In [1], Berg, Jackson and Jordán proved an interesting edge-splitting theorem for directed hypergraphs, which led to a solution for the problem of directed hypergraph edge-connectivity augmentation by uniform hyperarcs. In this note, we show that their edge-splitting result can be formulated in a more general form (using essentially the same proof). The result gives a method for solving a broader class of undirected and directed augmentation problems where the new hyperedges have the same prescribed size. In Section 3 we study problems where the aim is to obtain a directed hypergraph that covers a given crossing supermodular set function; this includes the problem of (k, l) -edge-connectivity augmentation. In Section 4 the objective is to obtain an undirected hypergraph that has an orientation covering a non-negative crossing supermodular set function. A notable special case is the (k, l) -partition-connectivity augmentation of undirected hypergraphs.

Let V be a finite ground set. For a function $m : V \rightarrow \mathbb{R}$ and a set $X \subseteq V$, we use the notation $m(X) := \sum_{v \in X} m(v)$. Hyperedges are considered to be multisets, so a *hyperedge* can be defined as a function $e : V \rightarrow \mathbb{Z}_+$, but we use the notations $v \in e$ for $e(v) > 0$ and $|e \cap X|$ for $e(X)$. A *hyperarc* a is a hyperedge with a designated *head node* $h(a) \in a$; the rest of its nodes are called *tail nodes* ($t(a) = a - h(a)$). An *orientation* of a hypergraph $H = (V, \mathcal{E})$ is a directed hypergraph obtained by

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designating a head node $h(e)$ for every $e \in \mathcal{E}$. A ν -hyperedge is a hyperedge e with $|e| = \nu$, while an $(r, 1)$ -hyperarc is a hyperarc a with $|t(a)| = r$. A hyperedge e enters a set X if $e \cap X \neq \emptyset$ and $e \cap (V - X) \neq \emptyset$, while a hyperarc a enters X if $h(a) \in X$ and $t(a) \cap (V - X) \neq \emptyset$. For a hypergraph $H = (V, \mathcal{E})$ and a directed hypergraph $D = (V, \mathcal{A})$ we define $d_H(X) := |\{e \in \mathcal{E} \mid e \text{ enters } X\}|$, $\varrho_D(X) := |\{a \in \mathcal{A} \mid a \text{ enters } X\}|$ and $\delta_D(X) = \varrho_D(V - X)$, which have the following properties:

$$d_H(X) + d_H(Y) \geq d_H(X \cap Y) + d_H(X \cup Y) \quad \text{for every } X, Y \subseteq V, \quad (1)$$

$$\varrho_D(X) + \varrho_D(Y) \geq \varrho_D(X \cap Y) + \varrho_D(X \cup Y) \quad \text{for every } X, Y \subseteq V, \quad (2)$$

$$\delta_D(X) + \delta_D(Y) \geq \delta_D(X \cap Y) + \delta_D(X \cup Y) \quad \text{for every } X, Y \subseteq V. \quad (3)$$

For a family \mathcal{F} of subsets of V , we use the notation $\text{co}(\mathcal{F}) := \{V - X \mid X \in \mathcal{F}\}$. Two sets X and Y are crossing if all of $X - Y, Y - X, X \cap Y, V - (X \cup Y)$ are non-empty. A family of sets is *cross-free* if it does not contain two crossing members. Let $p : 2^V \rightarrow \mathbb{Z} \cup \{-\infty\}$ be a set function (we always assume that $p(\emptyset) = 0$). A hypergraph H (a directed hypergraph D) is said to *cover* p if $d_H(X) \geq p(X)$ ($\varrho_D(X) \geq p(X)$) for every $X \subseteq V$. The set function p is *crossing supermodular* if

$$p(X) + p(Y) \leq p(X \cap Y) + p(X \cup Y) \quad (4)$$

holds whenever $X \cap Y \neq \emptyset$ and $V - (X \cup Y) \neq \emptyset$.

2 Directed splitting off

A special case of the following theorem (when $p(X) = k$ for every $\emptyset \neq X \subset V$ for some positive integer k) was proved in [1]. Here we show that a more general result can be proved with essentially the same techniques.

Theorem 2.1. *Let $p : 2^V \rightarrow \mathbb{Z} \cup \{-\infty\}$ be a crossing supermodular set function, $m_i : V \rightarrow \mathbb{Z}_+$ and $m_o : V \rightarrow \mathbb{Z}_+$ degree specifications such that $m_o(V) = r m_i(V)$ for some positive integer r , and*

$$m_i(X) \geq p(X) \quad \text{for every } X \subseteq V, \quad (5)$$

$$m_o(V - X) \geq p(X) \quad \text{for every } X \subseteq V. \quad (6)$$

Then there is a directed $(r, 1)$ -hypergraph D such that $\delta_D(v) = m_o(v)$ and $\varrho_D(v) = m_i(v)$ for every $v \in V$, and

$$\varrho_D(X) \geq p(X) \quad \text{for every } X \subseteq V.$$

Proof. Consider a hyperarc a for which $m_i(h(a)) > 0$, and $m_o(v) \geq |t(a) \cap \{v\}|$ for every $v \in t(a)$. We define vectors $m_i^a : V \rightarrow \mathbb{Z}_+$, $m_o^a : V \rightarrow \mathbb{Z}_+$, and a set function $p^a : 2^V \rightarrow \mathbb{Z} \cup \{-\infty\}$ the following way: m_i^a is obtained from m_i by decreasing it by 1 on $h(a)$, m_o^a is obtained from m_o by decreasing it on the nodes of $t(a)$ by their multiplicities in $t(a)$, and p^a is obtained by decreasing p by 1 on every set entered by a . The hyperarc a can be *split off* if $m_i^a(X) \geq p^a(X)$ and $m_o^a(V - X) \geq p^a(X)$ for every $X \subseteq V$. The operation is called a *feasible $(l, 1)$ -splitting* if $|t(a)| = l$. Note that p^a is crossing supermodular by (2). The following lemma describes conditions when a feasible splitting is available.

Lemma 2.2. *Let $p : 2^V \rightarrow \mathbb{Z} \cup \{-\infty\}$ be a crossing supermodular set function, $m_i : V \rightarrow \mathbb{Z}_+$ and $m_o : V \rightarrow \mathbb{Z}_+$ degree specifications such that $m_i(V) \leq m_o(V) \leq rm_i(V)$ for some integer r , and*

$$m_i(X) \geq p(X) \text{ for every } X \subseteq V, \quad (7)$$

$$m_o(V - X) \geq p(X) \text{ for every } X \subseteq V. \quad (8)$$

Let $u \in V$ be such that $m_i(u) > 0$. Then there is a hyperarc a with $h(a) = u$ and $|t(a)| \leq r$ that can be split off.

Proof. We can assume that $m_i(V) \geq 2$. A set X is called *in-critical* if $u \in X$ and $p(X) = m_i(X)$. The maximal in-critical sets are pairwise co-disjoint, since they intersect, and by the crossing supermodularity of p , the union of two crossing in-critical sets is in-critical. The complement of a maximal in-critical set is called a *petal*. Let \mathcal{F} denote the family of maximal in-critical sets, and let $\alpha := |\mathcal{F}|$; \mathcal{F} is called an α -*flower*.

Claim 2.3. $\alpha \leq r$.

Proof. Otherwise we would have

$$\sum_{X \in \mathcal{F}} p(X) = \sum_{X \in \mathcal{F}} m_i(X) > rm_i(V) \geq m_o(V) \geq \sum_{X \in \mathcal{F}} m_o(V - X),$$

which contradicts (8). □

First, suppose that $\alpha = 1$ ($a = \{u\}$ is obviously good for $\alpha = 0$), and let P be the single petal; $m_o(P) \geq m_i(V - P) > 0$. A set X is called *out-critical* if $u \notin X$ and $m_o(V - X) = p(X) > 0$; if there are no such sets, then for any $v \in P$ with $m_o(v) > 0$ the arc $a = vu$ can be split off. By the crossing supermodularity of p , the non-empty intersection of two out-critical sets is also out-critical. Since $u \notin X$ and $m_i(X) \geq m_o(V - X)$ for any out-critical set, there are no two disjoint out-critical sets, so there is a unique minimal out-critical Y . One of $P - Y$ and $Y - P$ is empty, otherwise $m_o(P - Y) + m_i(Y - P) < m_o(V - Y) + m_i(V - P) = p(Y) + p(V - P) \leq p(Y - P) + p(V - (P - Y)) \leq m_i(Y - P) + m_o(P - Y)$ would be a contradiction. Also, $m_o(Y) = m_o(V) - m_o(V - Y) \geq m_i(V) - m_o(V - Y) \geq m_i(V) - m_i(Y) > 0$, hence $m_o(Y \cap P) > 0$. Let $v \in Y \cap P$ be a node with $m_o(v) > 0$. Then the arc $a = vu$ can be split off.

If $\alpha \geq 2$, we define a by selecting as tail nodes one arbitrary node v with $m_o(v) > 0$ from each petal. We prove that a can be split off. By the construction of a , (7) holds after the splitting.

Suppose that there is a set X which violates (8) after the splitting, i.e. $m_o^a(V - X) < p^a(X)$. This means that if a enters X , then $p(X) > m_o(V - X) - |(t(a)) \cap (V - X)| + 1$ and if a does not enter X , then $u \notin X$ and $p(X) > m_o(V - X) - |(t(a)) \cap (V - X)|$. In both cases $p(X) > m_o(V - X) - |a \cap (V - X)| + 1$.

There is a petal P such that $P - X \neq \emptyset$ and $X - P \neq \emptyset$ (this is trivial if X is subset of a petal; if it is not, then any petal P is good for which $P \cap a \notin X$, and such a petal

exists otherwise $m_o(V - X) = m_o^a(V - X) < p^a(X) \leq p(X)$ contradicting (8)). The crossing supermodularity of p implies that

$$m_i(V - P) + m_o(V - X) - |a \cap (V - X)| + 1 < p(V - P) + p(X) \leq p(X - P) + p(V - (P - X)),$$

so

$$\begin{aligned} m_i(X - P) + m_o(P - X) - |a \cap (P - X)| + 1 &< p(X - P) + p(V - (P - X)), \\ m_o(P - X) - |a \cap (P - X)| + 1 &< p(V - (P - X)), \end{aligned}$$

which would imply that $V - (P - X)$ violates (8), since $|a \cap (P - X)| \leq 1$. \square

Proof of Theorem 2.1: According to Lemma 2.2 we can obtain a directed hypergraph D^* by successive feasible splitting off operations such that $\delta_{D^*}(v) = m_o(v)$, $\varrho_{D^*}(v) = m_i(v)$ for every $v \in V$, and $\varrho_{D^*}(X) \geq p(X)$ for every $X \subseteq V$. Since $m_o(V) = rm_i(V)$, $|a| \geq r + 1$ holds for at least one hyperarc a of D^* . So there is a feasible $(r_1, 1)$ -splitting with head $h(a)$ for some $r_1 \geq r$; moreover, by Lemma 2.2 there is also a feasible $(r_2, 1)$ -splitting with head $h(a)$ for some $r_2 \leq r$.

Lemma 2.4. *If for some $r_1 > r > r_2$ there is a feasible $(r_1, 1)$ -splitting and a feasible $(r_2, 1)$ -splitting with head u , then there is a feasible $(r, 1)$ -splitting with head u .*

Proof. Let a be the hyperarc obtained by the $(r_1, 1)$ -splitting. By induction, it suffices to show that for some $v \in t(a)$, the hyperarc a' defined by $h(a') = u$, $t(a') = t(a) - v$ gives a feasible splitting. If $a(v) \geq 2$ for some v we are ready, so suppose $a \leq 1$. If $r_1 > 2$, then suppose indirectly that for every $v \in t(a)$ there is an in-critical set X_v such that $a - v \subseteq X_v$ and $v \notin X_v$. We can assume that these are maximal in-critical sets. Thus the sets $\{X_v \mid v \in t(a)\}$ form a flower with r_1 petals centered on u ; but this contradicts the fact that there is a feasible $(r_2, 1)$ -splitting with head u .

If $r_1 = 2$, then $r_2 = 0$, so there are no in-critical sets. As we have seen, there is a unique minimal out-critical set Y with $u \notin Y$. Then $t(a) - Y = \emptyset$, otherwise the $(2, 1)$ -splitting would not be feasible; thus both $(1, 1)$ -splittings are feasible. \square

We prove Theorem 2.1 by induction on $m_i(V)$. According to Lemma 2.4, there is a node $u \in V$ with $m_i(u) > 0$ for which there exists a feasible $(r, 1)$ -splitting at u ; let a be the resulting $(r, 1)$ -hyperarc. By induction, there is a directed $(r, 1)$ -hypergraph D' that satisfies the conditions given by m_i^a , m_o^a and p^a . The directed hypergraph obtained by adding a to D' satisfies the conditions of Theorem 2.1. \square

3 Directed hypergraph augmentation

As in [1], one can obtain an edge-connectivity augmentation result from Theorem 2.1, using the following theorem of Fujishige:

Theorem 3.1 (Fujishige). *Let $p : 2^V \rightarrow \mathbb{Z} \cup \{-\infty\}$ be a crossing supermodular function and let*

$$B(p) := \{x \in \mathbb{R}^V \mid x(Z) \geq p(Z) \forall Z \subseteq V; x(V) = p(V)\}. \quad (9)$$

Then $B(p)$ is nonempty if and only if

$$\sum_{i=1}^t p(X_i) \leq p(V), \quad \sum_{i=1}^t p(V - X_i) \leq (t-1)p(V)$$

both hold for every partition $\{X_1, X_2, \dots, X_t\}$ of V . Furthermore, if $B(p)$ is nonempty, then it is a base polyhedron, thus its vertices are integral. \square

Theorem 3.2. Let $p : 2^V \rightarrow \mathbb{Z} \cup \{-\infty\}$ be a crossing supermodular set function. There exists a directed $(r, 1)$ -hypergraph with γ hyperarcs that covers p if and only if

$$\gamma \geq \sum_{X \in \mathcal{F}} p(X), \quad (10)$$

$$r\gamma \geq \sum_{X \in \mathcal{F}} p(V - X), \quad (11)$$

$$(|\mathcal{G}| - 1)\gamma \geq \sum_{X \in \mathcal{G}} p(V - X) \quad (12)$$

hold for every sub-partition \mathcal{F} and for every partition \mathcal{G} of V .

Proof. The necessity of the conditions can be seen easily. To prove sufficiency, one can construct degree specifications m_i and m_o that satisfy the conditions of Theorem 2.1. Let us define the set function $p' : 2^V \rightarrow \mathbb{Z}_+$ by $p'(V) = \gamma$, $p'(x) = \max\{0, p(x)\}$ for singletons and $p'(X) = p(X)$ otherwise. Note that p' is crossing supermodular. If \mathcal{F} is a partition of V , then (10) implies that

$$\sum_{X \in \mathcal{F}} p'(X) \leq \sum_{X \in \mathcal{F}, p(X) > 0} p(X) \leq \gamma,$$

and either (12), or (10) applied to sub-partitions with one class, implies that

$$\sum_{X \in \text{co}(\mathcal{F})} p'(X) \leq \sum_{X \in \text{co}(\mathcal{F}), p(X) > 0} p(X) \leq (|\mathcal{F}| - 1)\gamma.$$

Thus, by applying Theorem 3.1 to p' , we get a nonnegative integer vector m_i s.t. $m_i(X) \geq p(X)$ for all $X \subseteq V$, and $m_i(V) = \gamma$.

To construct m_o consider a nonnegative vector satisfying $m_o(V - X) \geq p(X)$ for all $X \subseteq V$, which is minimal in the sense that for every $v \in V$ with $m_o(v) > 0$, there exists a set X for which $v \notin X$ and $m_o(V - X) = p(X)$. Choose a family $\mathcal{G} = \{X_1, X_2, \dots, X_l\}$ of such sets with l minimal. Hence, two sets cannot cross, since we could replace them by their intersection. If the family is composed of co-disjoint sets, then

$$m_o(V) = \sum_{i=1}^l m_o(V - X_i) = \sum_{i=1}^l p(X_i) \leq r\gamma.$$

by (11). If there are two disjoint sets, X_i and X_j , then

$$m_o(V) \leq m_o(V - X_i) + m_o(V - X_j) = p(X_i) + p(X_j) \leq \gamma.$$

Now we can increase m_o on an arbitrary node to obtain $m_o(V) = r\gamma$, and apply Theorem 2.1 to construct a directed $(r, 1)$ -hypergraph with degrees m_i and m_o that covers p . \square

The following example demonstrates that condition (12) cannot be left out. Let $V = \{v_1, v_2, v_3\}$, $p(\{v_1, v_2\}) = p(\{v_1, v_3\}) = p(\{v_2, v_3\}) = 2$ and $p(X) = 0$ for the other sets, $r = 3$, $\gamma = 2$. Conditions (10) and (11) are satisfied, but (12) is not, and there is no directed $(3, 1)$ -hypergraph of 2 hyperarcs covering p . In the graph case (12) easily follows from (11).

A special case where (12) follows from (10) is the (k, l) -edge-connectivity augmentation of directed hypergraphs (for $l \leq k$), which is a generalization of the k -edge-connectivity augmentation problem studied in [1]. For a directed hypergraph D and a fixed $s \in V$, let $p(X) = k - \varrho_D(X)$ if $s \notin X \neq \emptyset$, $p(X) = l - \varrho_D(X)$ if $s \in X \neq V$ and $p(X) = 0$ otherwise. Let $\mathcal{F} = \{X_1, X_2, \dots, X_t\}$ be a partition of V ($t \geq 2$). If (10) holds, then $\sum_{X \in \text{co}(\mathcal{F})} p(X) = (t-1)l + k - \sum_{X \in \text{co}(\mathcal{F})} \varrho_D(X) \leq l + (t-1)k - \sum_{X \in \text{co}(\mathcal{F})} \delta_D(X) = l + (t-1)k - \sum_{X \in \mathcal{F}} \varrho_D(X) = \sum_{X \in \mathcal{F}} p(X) \leq \gamma$.

Corollary 3.3. *A directed hypergraph $D = (V, \mathcal{A})$ can be made (k, l) -edge-connected with γ new $(r, 1)$ -hyperarcs if and only if*

$$\begin{aligned} \gamma &\geq \sum_{X \in \mathcal{F}} p(X), \\ r\gamma &\geq \sum_{X \in \mathcal{F}} p(V - X) \end{aligned}$$

hold for every subpartition \mathcal{F} of V , where p is the above-defined set-function.

4 Augmentation and orientation

In this section we consider only non-negative crossing supermodular set functions. We are interested in the problem of adding ν -hyperedges to an initial undirected hypergraph, so that the resulting hypergraph has an orientation covering a given set function p . Similar problems for graphs were studied in [2]. As in that case, we first solve the degree specified problem, and then obtain a min-max formula for minimum cardinality augmentation. Some new notations are introduced to facilitate the formulation of the min-max results.

A family \mathcal{F} of sets is a *composition* of a set $X \subseteq V$ if the value $\sum_{Z \in \mathcal{F}} \chi_Z(v) - \chi_X(v)$ is the same for every $v \in V$. A composition of V is called a *regular family*. For a set X and a family \mathcal{F} that is a composition of X , let

$$\alpha_X(\mathcal{F}) := \sum_{Z \in \mathcal{F}} \chi_Z(v) - \chi_X(v) \quad \text{for an arbitrary } v \in V.$$

A composition of X is a *tree-composition* if it is cross-free and it contains no proper subfamily that is a partition or a co-partition of V . Tree-compositions have the following properties:

Claim 4.1. *If $\mathcal{F} \neq \emptyset$ is a tree-composition of X that is not a partition of X , then it contains a subfamily $\{Z_1, \dots, Z_t\}$ ($t \geq 2$) of pairwise co-disjoint sets such that $\cap Z_i \subseteq X$. If $X \neq V$, then $Z_i - X \neq \emptyset$ ($i = 1, \dots, t$). \square*

For a hypergraph $H = (V, \mathcal{E})$ and a set $X \subseteq V$, let $i_H(X)$ denote the number of hyperedges $e \in \mathcal{E}$ with $e \cap (V - X) = \emptyset$. For a regular family \mathcal{F} let

$$e_H(\mathcal{F}) := \alpha_\emptyset(\mathcal{F})|\mathcal{E}| - \sum_{X \in \mathcal{F}} i_H(X). \quad (13)$$

More intuitively, $e_H(\mathcal{F}) = \max\{\sum_{X \in \mathcal{F}} \varrho_{\vec{H}}(X) \mid \vec{H} \text{ is an orientation of } H\}$.

Theorem 4.2. *Let $H = (V, \mathcal{E})$ be a hypergraph, $p : 2^V \rightarrow \mathbb{Z}_+$ a non-negative crossing supermodular set function, and $m : V \rightarrow \mathbb{Z}_+$ a degree specification where $m(V)$ is divisible by a fixed integer $\nu \geq 2$. There exists a ν -uniform hypergraph I such that $H + I$ has an orientation covering p and $d_I(v) = m(v)$ for every $v \in V$ if and only if the following hold for every partition \mathcal{F} of V :*

$$\frac{m(V)}{\nu} \geq \sum_{Z \in \mathcal{F}} p(Z) - e_H(\mathcal{F}), \quad (14)$$

$$\min_{X \in \mathcal{F}} m(V - X) \geq \sum_{Z \in \mathcal{F}} p(Z) - e_H(\mathcal{F}), \quad (15)$$

$$\min_{\mathcal{F}' \subseteq \mathcal{F}, X = \cup \mathcal{F}'} \left(m(V - X) + (|\mathcal{F}'| - 1) \frac{m(V)}{\nu} \right) \geq \sum_{Z \in \mathcal{F}} p(V - Z) - e_H(\text{co}(\mathcal{F})). \quad (16)$$

Proof. The right hand side of the inequalities is the deficiency of the hyperedges of H . The necessity of the conditions follows from the observation that the left hand side is always an upper bound on the contribution of the new hyperarcs. In (14): every new hyperarc can enter at most one set of \mathcal{F} ; in (15): every hyperarc that enters a set of \mathcal{F} must have a node in $V - X$; in (16): the number of sets of $\text{co}(\mathcal{F})$ that a new hyperarc enters is at most $|\mathcal{F}'| - 1$ plus the number of nodes it has in $V - X$.

The proof of sufficiency is similar to that of the augmentation theorem in [2]. We add a new node z to the set of nodes, and for every $v \in V$ we add $m(v)$ parallel edges between v and z ; the resulting hypergraph is denoted by $H' = (V', \mathcal{E}')$. Our first aim is to find an orientation \vec{H}' of H' that has the following properties:

$$\varrho_{\vec{H}'}(V) = \frac{m(V)}{\nu}, \quad (17)$$

$$\varrho_{\vec{H}'}(X) \geq p(X) \quad \text{if } \emptyset \neq X \subset V, \quad (18)$$

$$\varrho_{\vec{H}'}(X + z) \geq p(X) \quad \text{if } \emptyset \neq X \subset V. \quad (19)$$

To find such an orientation, we use the following lemma (see e.g. [3]):

Lemma 4.3. *Given a hypergraph H' and a vector $x' : V' \rightarrow \mathbb{Z}_+$, there is an orientation \vec{H}' of H' such that $\varrho_{\vec{H}'}(v) = x'(v)$ for every $v \in V'$ if and only if $x'(V') = |\mathcal{E}'|$ and $x'(Y) \geq i_{H'}(Y)$ for every $Y \subseteq V'$. \square*

We call a vector $x : V \rightarrow \mathbb{Z}_+$ *feasible* if it is the vector of in-degrees (restricted to V) of an orientation satisfying (17)–(19). It is easy to see using Lemma 4.3 that x is feasible if and only if $x(V) = |\mathcal{E}| + \frac{m(V)}{\nu}$ and $x(Z) \geq p_m(Z)$ for every $Z \subseteq V$, where

$$p_m(X) := p(X) + i_H(X) + \left(m(X) - \frac{\nu-1}{\nu} m(V) \right)^+ \quad (X \subseteq V). \quad (20)$$

The set function p_m is crossing supermodular. A vector x is feasible if and only if it is an integral element of $B(p_m)$ (as defined in (9)).

Claim 4.4. *If conditions (14)–(16) are satisfied, then $B(p_m)$ is non-empty.*

Proof. By Theorem 3.1, it suffices to show that

$$\sum_{X \in \mathcal{F}} p_m(X) \leq |\mathcal{E}| + \frac{m(V)}{\nu}, \quad (21)$$

$$\sum_{X \in \mathcal{F}} p_m(V - X) \leq (|\mathcal{F}| - 1) \left(|\mathcal{E}| + \frac{m(V)}{\nu} \right) \quad (22)$$

for every partition \mathcal{F} . Note that $m(X) - \frac{\nu-1}{\nu} m(V)$ can be positive for at most one member of a partition. Thus (21) follows from (13), and either (15) or (14), depending on whether \mathcal{F} has such a member or not. The inequality (22) follows from (13) and (16). \square

By Theorem 3.1, $B(p_m)$ is a base polyhedron with integral vertices, and any such vertex x is the vector of in-degrees (restricted to V) of an orientation \vec{H}' satisfying (17)–(19).

Let $m_i(v)$ be the multiplicity of the arc zv in \vec{H}' , $m_o(v)$ be the multiplicity of the arc vz in \vec{H}' , and let \vec{H} denote the directed hypergraph obtained from \vec{H}' by deleting the node z . Then $m_i(X) \geq p(X) - \varrho_{\vec{H}}(X)$ and $m_o(V - X) \geq p(X) - \varrho_{\vec{H}}(X)$ for every $X \subseteq V$. By (2) and the crossing supermodularity of p , the set function $q(X) := p(X) - \varrho_{\vec{H}}(X)$ is crossing supermodular. Theorem 2.1 asserts the existence of a directed $(\nu-1, 1)$ -hypergraph D that covers q , and satisfies the degree specifications m_i and m_o . This means that $\vec{H} + D$ covers p , and the undirected hypergraph I that underlies D satisfies the degree specification m . Since $\vec{H} + D$ is an orientation of $H + I$, this completes the proof of Theorem 4.2. \square

Given a characterization of the degree specifications that allow a good augmentation, it is often possible to deduce a characterization of the minimum number of hyperedges needed. In the present case we obtain the following theorem:

Theorem 4.5. *Let $H = (V, \mathcal{E})$ be a hypergraph, $p : 2^V \rightarrow \mathbb{Z}_+$ a non-negative crossing supermodular set function, and $\nu \geq 2$ an integer. There exists a ν -uniform hypergraph I with γ hyperedges such that $H + I$ has an orientation covering p if and only if*

$$\gamma(\nu + \alpha_X(\mathcal{F}_1) + (\nu - 1)\alpha_X(\mathcal{F}_2)) \geq \sum_{Z \in \mathcal{F}_1 + \text{co}(\mathcal{F}_2)} p(Z) - e_H(\mathcal{F}_1 + \text{co}(\mathcal{F}_2)) \quad (23)$$

whenever \mathcal{F}_1 and \mathcal{F}_2 are tree-compositions of some set $X \subseteq V$, $\mathcal{F}_1 + \mathcal{F}_2$ is cross-free, and $\alpha_X(\mathcal{F}_2) \leq 0$ (i.e. either \mathcal{F}_2 is a partition of X , or $X = V$ and $\mathcal{F}_2 = \emptyset$).

Proof. The right hand side of (23) is the deficiency of H relative to the family $\mathcal{F}_1 + \text{co}(\mathcal{F}_2)$. The number of sets of \mathcal{F}_1 that a new hyperarc enters is at most $\alpha_X(\mathcal{F}_1)$, plus 1 if its head is in X . The number of sets of $\text{co}(\mathcal{F}_2)$ that a new hyperarc enters is at most $(\nu - 1)\alpha_X(\mathcal{F}_2)$ plus the number of tail nodes it has in X . This shows the necessity of (23). To prove sufficiency, we define for every $X \subseteq V$ and compositions $\mathcal{F}_1, \mathcal{F}_2$ of X :

$$Q_X(\mathcal{F}_1, \mathcal{F}_2) := \sum_{Z \in \mathcal{F}_1 + \text{co}(\mathcal{F}_2)} p(Z) - e_H(\mathcal{F}_1 + \text{co}(\mathcal{F}_2)) - \gamma(\alpha_X(\mathcal{F}_1) + (\nu - 1)\alpha_X(\mathcal{F}_2)),$$

$$q(X) := \max\{Q_X(\mathcal{F}_1, \mathcal{F}_2) : \mathcal{F}_1 \text{ and } \mathcal{F}_2 \text{ are tree-compositions of } X, \\ \mathcal{F}_1 + \mathcal{F}_2 \text{ is cross-free, } \alpha_X(\mathcal{F}_2) \leq 0\}.$$

Condition (23) is equivalent to the inequality $\max_{X \subseteq V} q(X) \leq \nu\gamma$; let us assume that this holds. We can observe that if $m : V \rightarrow \mathbb{Z}_+$ satisfies $m(X') \geq q(X')$ for every $X' \subseteq V$ and $m(V) = \nu\gamma$, then m satisfies (14)–(16). The choice where $X' = V$, $\mathcal{F}_1 = \mathcal{F}$ is a partition of V , $\mathcal{F}_2 = \emptyset$, $(\alpha_{X'}(\mathcal{F}_1) = 0, \alpha_{X'}(\mathcal{F}_2) = -1)$ easily yields (14), by $\gamma = \frac{m(V)}{\nu}$. With $X' = V - X$, where $\mathcal{F}_1 = \mathcal{F} - \{X\}$ is a partition of X' and $\mathcal{F}_2 = \{X\}$, $(\alpha_{X'}(\mathcal{F}_1) = 0, \alpha_{X'}(\mathcal{F}_2) = 0)$, (15) follows. To obtain (16), we set $X' = V - X$, $\mathcal{F}_1 = \text{co}(\mathcal{F}')$, $\mathcal{F}_2 = \mathcal{F} - \mathcal{F}'$, $(\alpha_{X'}(\mathcal{F}_1) = |\mathcal{F}_1| - 1, \alpha_{X'}(\mathcal{F}_2) = 0)$. Thus by Theorem 4.2 the existence of such an m implies the existence of a hypergraph I that satisfies the requirements. To prove that such an m exists, we use the properties of a set function slightly different from q :

$$q'(X) := \max\{Q_X(\mathcal{F}_1, \mathcal{F}_2) : \mathcal{F}_1 \text{ and } \mathcal{F}_2 \text{ are tree-compositions of } X\}.$$

Claim 4.6. *The value $Q_X(\mathcal{F}_1, \mathcal{F}_2)$ does not decrease if we remove a partition or a co-partition of V from \mathcal{F}_1 or \mathcal{F}_2 .*

Proof. It is easy to see that if $X \cap Y = \emptyset$, $\mathcal{F}_1^X, \mathcal{F}_2^X$ are compositions of X , and $\mathcal{F}_1^Y, \mathcal{F}_2^Y$ are compositions of Y , then

$$Q_X(\mathcal{F}_1^X, \mathcal{F}_2^X) + Q_Y(\mathcal{F}_1^Y, \mathcal{F}_2^Y) = Q_{X \cup Y}(\mathcal{F}_1^X + \mathcal{F}_1^Y, \mathcal{F}_2^X + \mathcal{F}_2^Y). \quad (24)$$

The case $Y = \emptyset$ proves the claim, since $q(V) \leq \nu\gamma$ implies that $q(\emptyset) \leq 0$. \square

Claim 4.7. *The set function q' is fully supermodular.*

Proof. Let $X, Y \subseteq V$, and suppose that the maximum in the definition of q' is reached on families $\mathcal{F}_1^X, \mathcal{F}_2^X$, and $\mathcal{F}_1^Y, \mathcal{F}_2^Y$, respectively. Let $\mathcal{F}_1 := \mathcal{F}_1^X + \mathcal{F}_1^Y$, $\mathcal{F}_2 := \mathcal{F}_2^X + \mathcal{F}_2^Y$. We apply the following operations, as long as any of them is possible:

- If $Z_1, Z_2 \in \mathcal{F}_1$ are crossing, then replace them in \mathcal{F}_1 by $Z_1 \cap Z_2$ and $Z_1 \cup Z_2$.
- If $Z_1, Z_2 \in \mathcal{F}_2$ are crossing, then replace them in \mathcal{F}_2 by $Z_1 \cap Z_2$ and $Z_1 \cup Z_2$.

It is easy to see that after a finite number of steps, the resulting families \mathcal{F}'_1 and \mathcal{F}'_2 become cross-free. Then \mathcal{F}'_i decomposes into a composition $\mathcal{F}'_i^{X \cap Y}$ of $X \cap Y$ and a composition $\mathcal{F}'_i^{X \cup Y}$ of $X \cup Y$ ($i = 1, 2$); and all of these families are cross-free. The crossing supermodularity of p implies that $\sum_{Z \in \mathcal{F}'_1^X + \mathcal{F}'_1^Y} p(Z) \leq \sum_{Z \in \mathcal{F}'_1^{X \cap Y} + \mathcal{F}'_1^{X \cup Y}} p(Z)$ and $\sum_{Z \in \mathcal{F}'_2^X + \mathcal{F}'_2^Y} p(V - Z) \leq \sum_{Z \in \mathcal{F}'_2^{X \cap Y} + \mathcal{F}'_2^{X \cup Y}} p(V - Z)$. It is easy to check that $e_H(\mathcal{F}'_1^X + \text{co}(\mathcal{F}'_2^X)) + e_H(\mathcal{F}'_1^Y + \text{co}(\mathcal{F}'_2^Y)) \geq e_H(\mathcal{F}'_1^{X \cap Y} + \text{co}(\mathcal{F}'_2^{X \cap Y})) + e_H(\mathcal{F}'_1^{X \cup Y} + \text{co}(\mathcal{F}'_2^{X \cup Y}))$, and that $\alpha_X(\mathcal{F}'_1^X) + \alpha_Y(\mathcal{F}'_1^Y) = \alpha_{X \cap Y}(\mathcal{F}'_1^{X \cap Y}) + \alpha_{X \cup Y}(\mathcal{F}'_1^{X \cup Y})$ ($i = 1, 2$). Hence $Q(\mathcal{F}'_1^X, \mathcal{F}'_2^X) + Q(\mathcal{F}'_1^Y, \mathcal{F}'_2^Y) \leq Q(\mathcal{F}'_1^{X \cap Y}, \mathcal{F}'_2^{X \cap Y}) + Q(\mathcal{F}'_1^{X \cup Y}, \mathcal{F}'_2^{X \cup Y})$; using Claim 4.6, we obtain that $q'(X) + q'(Y) \leq q'(X \cap Y) + q'(X \cup Y)$. \square

Claim 4.7 and Theorem 3.1 imply that there exists a vector $m : V \rightarrow \mathbb{Z}_+$ with $m(V) = \nu\gamma$ that satisfies $m(X) \geq q'(X)$ for every $X \subseteq V$ if and only if $\max_{X \subseteq V} q'(X) \leq \nu\gamma$.

Claim 4.8. *If condition (23) holds, then $\max_{X \subseteq V} q'(X) = \max_{X \subseteq V} q(X) \leq \nu\gamma$.*

Proof. Let X be the set where the maximum is reached for q' , and let $\mathcal{F}_1, \mathcal{F}_2$ be tree-compositions of X for which $q'(X) = Q_X(\mathcal{F}_1, \mathcal{F}_2)$. We transform \mathcal{F}_1 and \mathcal{F}_2 using the following operations until none of them is applicable:

- If $Z_1, Z_2 \in \mathcal{F}_1$ are crossing, then replace Z_1, Z_2 by $Z_1 \cap Z_2, Z_1 \cup Z_2$ in \mathcal{F}_1 .
- If $Z_1, Z_2 \in \mathcal{F}_2$ are crossing, then replace Z_1, Z_2 by $Z_1 \cap Z_2, Z_1 \cup Z_2$ in \mathcal{F}_2 .
- If \mathcal{F}_2 is a partition of some $Z \subseteq V$, and $Z_1 \in \mathcal{F}_1$ and $Z_2 \in \mathcal{F}_2$ are crossing, then replace Z_1 by $Z_1 - Z_2$ in \mathcal{F}_1 , and replace Z_2 by $Z_2 - Z_1$ in \mathcal{F}_2 .
- If $\{Z_1, \dots, Z_t\} \subset \mathcal{F}_1$ or $\{Z_1, \dots, Z_t\} \subset \mathcal{F}_2$ is a partition or a co-partition of V , then remove Z_1, \dots, Z_t from that family.
- If \mathcal{F}_2 is a composition of $Z \subseteq V$ and it contains a subfamily $\{Z_1, \dots, Z_t\}$ ($t \geq 2$) of pairwise co-disjoint sets such that $\emptyset \neq \cap Z_i \subseteq Z$, then remove Z_1, \dots, Z_t from \mathcal{F}_2 , and add $V - Z_1, \dots, V - Z_t$ to \mathcal{F}_1 .

It is easy to see that this terminates after a finite number of steps. We denote by \mathcal{F}'_1 and \mathcal{F}'_2 the families obtained at the end of the process. Then, by Claim 4.1, \mathcal{F}'_1 and \mathcal{F}'_2 are tree-compositions of some $X' \subseteq X$, $\alpha_{X'}(\mathcal{F}'_2) \leq 0$, and $\mathcal{F}'_1 + \mathcal{F}'_2$ is cross-free. Moreover, $Q_X(\mathcal{F}_1, \mathcal{F}_2)$ does not decrease in any of the steps (in the first 3 cases this follows from the supermodularity of p , in the 4th it is a consequence of Claim 4.6, and in the 5th it is obvious from the definition of $Q_X(\mathcal{F}_1, \mathcal{F}_2)$). This proves that $\max_{X \subseteq V} q'(X) = \max_{X \subseteq V} q(X)$. \square

By Claim 4.8, $\nu\gamma \geq \max_{X \subseteq V} q'(X)$. Thus there exists a vector $m : V \rightarrow \mathbb{Z}_+$ with $m(V) = \nu\gamma$ that satisfies (14)–(16), therefore by Theorem 4.2 there exists a ν -uniform hypergraph I with γ hyperedges such that $H + I$ has an orientation covering p . This concludes the proof of Theorem 4.5. \square

If the requirement function is monotone decreasing (i.e. $p(X) \geq p(Y)$ if $\emptyset \neq X \subseteq Y$), or symmetric, then the conditions of Theorem 4.2 and Theorem 4.5 can be simplified.

Theorem 4.9. *Let $H = (V, \mathcal{E})$ be a hypergraph, $p : 2^V \rightarrow \mathbb{Z}_+$ a monotone decreasing or symmetric non-negative crossing supermodular set function, and $m : V \rightarrow \mathbb{Z}_+$ a degree specification where $m(V)$ is divisible by a fixed integer $\nu \geq 2$. There exists a ν -uniform hypergraph I with degree-specification m such that $H + I$ has an orientation covering p if and only if the following hold for every partition \mathcal{F} of V :*

$$\frac{m(V)}{\nu} \geq \sum_{Z \in \mathcal{F}} p(Z) - e_H(\mathcal{F}), \quad (25)$$

$$\min_{X \in \mathcal{F}} m(V - X) \geq \sum_{Z \in \mathcal{F}} p(Z) - e_H(\mathcal{F}). \quad (26)$$

Proof. By definition, $e_H(\mathcal{F}) \leq e_H(\text{co}(\mathcal{F}))$ for every partition \mathcal{F} of V , and the monotonicity or symmetry of p implies that $\sum_{Z \in \text{co}(\mathcal{F})} p(Z) \leq \sum_{Z \in \mathcal{F}} p(Z)$ also holds. It is easy to see from this that (16) is implied by (14) if $|\mathcal{F}'| = 0$ or $|\mathcal{F}'| \geq 2$, and it is implied by (15) if $|\mathcal{F}'| = 1$. \square

Theorem 4.10. *Let $H = (V, \mathcal{E})$ be a hypergraph, $p : 2^V \rightarrow \mathbb{Z}_+$ a monotone decreasing or symmetric non-negative crossing supermodular set function, and $\nu \geq 2$ an integer. There exists a ν -uniform hypergraph I with γ hyperedges such that $H + I$ has an orientation covering p if and only if the following hold:*

$$\gamma \geq \sum_{Z \in \mathcal{F}} p(Z) - e_H(\mathcal{F}) \text{ for every partition } \mathcal{F}, \quad (27)$$

$$\nu\gamma \geq \sum_{Z \in \mathcal{F}_1 + \text{co}(\mathcal{F}_2)} p(Z) - e_H(\mathcal{F}_1 + \text{co}(\mathcal{F}_2)) \quad (28)$$

whenever \mathcal{F}_1 and \mathcal{F}_2 are partitions of some $X \subseteq V$ and \mathcal{F}_1 is a refinement of \mathcal{F}_2 .

Proof. It suffices to show that if condition (23) is violated for some pair $(\mathcal{F}_1, \mathcal{F}_2)$, then it is also violated by a pair $(\mathcal{F}'_1, \mathcal{F}'_2)$ that has the additional properties that \mathcal{F}'_1 is a partition of some $X' \subseteq V$, and $Y_2 \not\subseteq Y_1$ for every $Y_1 \in \mathcal{F}'_1, Y_2 \in \mathcal{F}'_2$. Such families can be obtained from \mathcal{F}_1 and \mathcal{F}_2 by repeating the following operations as long as any of them is possible:

- If \mathcal{F}_1 is a composition of $Z \subseteq V$, it contains a subfamily $\{W_1, \dots, W_s\}$ ($s \geq 2$) of pairwise co-disjoint sets such that $W := \cap W_i \subseteq Z$, and \mathcal{F}_2 contains a partition $\{Z_1, \dots, Z_t\}$ of W , then remove W_1, \dots, W_s from \mathcal{F}_1 and Z_1, \dots, Z_t from \mathcal{F}_2 .
- If \mathcal{F}_1 is a composition of $Z \subseteq V$ and it contains a set $W \subseteq Z$ such that \mathcal{F}_2 contains a partition $\{Z_1, \dots, Z_t\}$ of W , then replace W in \mathcal{F}_1 by the sets Z_1, \dots, Z_t , and replace Z_1, \dots, Z_t in \mathcal{F}_2 by W .

After a finite number of steps, none of the above operations are applicable; let $(\mathcal{F}'_1, \mathcal{F}'_2)$ be the pair obtained at that point. Then it follows from Claim 4.1 that \mathcal{F}'_1 must be a partition of some $X' \subseteq V$, and \mathcal{F}'_1 is a refinement of \mathcal{F}'_2 if $\mathcal{F}'_2 \neq \emptyset$. If

(27) holds, then the value $Q_X(\mathcal{F}_1, \mathcal{F}_2)$ does not decrease during the above two operations: in the first case this follows from (24), since $Q_W(\{W_1, \dots, W_s\}, \{Z_1, \dots, Z_t\}) \leq Q_\emptyset(\{V - W_1, \dots, V - W_s, Z_1, \dots, Z_t\}, \emptyset) \leq 0$; in the second case, it follows because $p(W) + \sum_{i=1}^t p(V - Z_i) \leq p(V - W) + \sum_{i=1}^t p(Z_i)$. \square

The above results have a straightforward application concerning the partition-connectivity augmentation of undirected hypergraphs. A hypergraph H is called (k, l) -*partition-connected* for non-negative integers $k \geq l$ if $e_H(\mathcal{F}) \geq (|\mathcal{F}| - 1)k + l$ for every partition \mathcal{F} . These hypergraphs have the following characterization:

Lemma 4.11 ([3]). *A hypergraph is (k, l) -partition-connected if and only if it has a (k, l) -edge-connected orientation.* \square

From Lemma 4.11 and Theorems 4.9 and 4.10 we obtain the following corollaries:

Corollary 4.12. *Let $H = (V, \mathcal{E})$ be a hypergraph, $m : V \rightarrow \mathbb{Z}_+$ a degree specification with $m(V)$ divisible by a fixed integer $\nu \geq 2$, and $k \geq l$ non-negative integers. There exists a ν -uniform hypergraph I such that $H + I$ is (k, l) -partition-connected and $d_H(v) = m(v)$ for all $v \in V$ if and only if the following hold for every partition \mathcal{F} of V :*

$$\frac{m(V)}{\nu} \geq (|\mathcal{F}| - 1)k + l - e_H(\mathcal{F}) , \quad (29)$$

$$\min_i m(V - X_i) \geq (|\mathcal{F}| - 1)k + l - e_H(\mathcal{F}) . \quad (30)$$

Corollary 4.13. *Let $H = (V, \mathcal{E})$ be a hypergraph, $\nu \geq 2$ and $k \geq l$ non-negative integers. There is a ν -uniform hypergraph I with γ edges such that $H + I$ is (k, l) -partition-connected if and only if the following two conditions are met:*

1. $\gamma \geq (|\mathcal{F}| - 1)k + l - e_H(\mathcal{F})$ for every partition \mathcal{F} ,
2. $\nu\gamma \geq |\mathcal{F}_1|k + |\mathcal{F}_2|l - e_H(\mathcal{F}_1 + \text{co}(\mathcal{F}_2))$ whenever \mathcal{F}_1 and \mathcal{F}_2 are partitions of some $X \subseteq V$ and \mathcal{F}_1 is a refinement of \mathcal{F}_2 .

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