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**Generically globally rigid zeolites  
in the plane**

Tibor Jordán

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# Generically globally rigid zeolites in the plane

Tibor Jordán\*

## Abstract

A  $d$ -dimensional zeolite is a  $d$ -dimensional body-and-pin framework with a  $(d+1)$ -regular underlying graph  $G$ . That is, each body of the zeolite is incident with  $d+1$  pins and each pin belongs to exactly two bodies. The corresponding  $d$ -dimensional combinatorial zeolite is a bar-and-joint framework whose graph is the line graph of  $G$ .

We show that a two-dimensional combinatorial zeolite is generically globally rigid if and only if its underlying 3-regular graph  $G$  is 3-edge-connected. The proof is based on a new rank formula for the two-dimensional rigidity matroid of line graphs.

## 1 Introduction

A  $d$ -dimensional zeolite is a  $d$ -dimensional body-and-pin framework in which each body is incident with  $d+1$  pins and each pin belongs to exactly two bodies. In the underlying graph  $G$  of the zeolite vertices correspond to bodies and two vertices are adjacent if and only if the corresponding bodies share a pin. Thus the underlying graph of the zeolite is  $(d+1)$ -regular.

By replacing the bodies by complete bar frameworks one obtains a  $d$ -dimensional combinatorial zeolite. It is a bar-and-joint framework whose graph is the line graph of the underlying graph  $G$  of the zeolite. (The *line graph*  $L(G)$  of a graph  $G = (V, E)$  is the simple graph with vertex set  $\{v_e : e \in E\}$ , where two vertices  $v_e, v_f$  are adjacent if and only if  $e, f$  have a common end-vertex in  $G$ .) See Figure 1 for a two-dimensional example.

The investigation of these structures is motivated in part by the existence (and flexibility properties) of real zeolites, which are molecules formed by corner-sharing tetrahedra, see e.g. [3]. Planar plate frameworks (which contain planar zeolites as a special case), in which each body is a regular polygon, have also been studied in the combinatorial rigidity literature [2]. In this paper we shall consider the (global) rigidity properties of planar combinatorial zeolites in generic position.

Roughly speaking, a combinatorial zeolite is globally rigid if its bar lengths uniquely determine the whole framework, up to congruence. Brigitte Servatius and Herman

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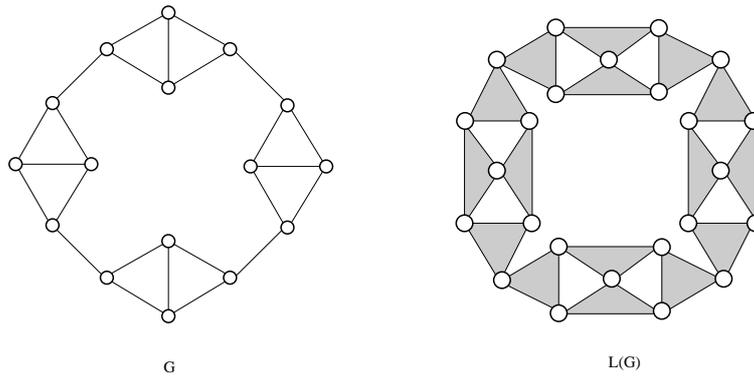


Figure 1: A 3-regular graph  $G$  and its line graph  $L(G)$ . The shaded triangles of the bar-and-joint framework on  $L(G)$  correspond to the bodies in the two-dimensional zeolite whose underlying graph is  $G$ .

Servatius [9] asked whether there is a simple necessary and sufficient condition, in terms of its underlying graph, for the global rigidity of a planar zeolite whose vertices are in generic position. We shall give an affirmative answer in Section 3 by showing that a planar combinatorial zeolite is generically globally rigid if and only if its 3-regular underlying graph is 3-edge-connected. The proof is based on a new rank formula for the two-dimensional rigidity matroid of line graphs. This formula, along with the necessary definitions, is given in Section 2. The last section is devoted to some concluding remarks.

## 2 Rigidity of line graphs

We shall need the following basic notions of combinatorial rigidity. For a detailed survey of the area we refer the reader to [1, 10]. A  $d$ -dimensional (bar-and-joint) *framework* is a pair  $(G, p)$ , where  $G = (V, E)$  is a graph and  $p$  is a map from  $V$  to  $\mathbb{R}^d$ . We also say that  $(G, p)$  is a *realization* of  $G$  in  $\mathbb{R}^d$ . We can think of the edges and vertices of  $G$  in the framework as rigid (fixed length) bars and universal joints, respectively. Two frameworks  $(G, p)$  and  $(G, q)$  are *equivalent* if corresponding edges have the same lengths, that is, if  $\|p(u) - p(v)\| = \|q(u) - q(v)\|$  holds for all pairs  $u, v$  with  $uv \in E$ , where  $\|\cdot\|$  denotes the Euclidean norm in  $\mathbb{R}^d$ . Frameworks  $(G, p)$ ,  $(G, q)$  are *congruent* if  $\|p(u) - p(v)\| = \|q(u) - q(v)\|$  holds for all pairs  $u, v$  with  $u, v \in V$ . We shall say that  $(G, p)$  is *globally rigid* if every framework which is equivalent to  $(G, p)$  is congruent to  $(G, p)$ .

Rigidity is a weaker property of frameworks than global rigidity. Intuitively, a framework is rigid if it has no continuous deformations. Equivalently, and more formally, a framework  $(G, p)$  is *rigid* if there exists an  $\epsilon > 0$  such that, if  $(G, q)$  is equivalent to  $(G, p)$  and  $\|p(u) - q(u)\| < \epsilon$  for all  $v \in V$ , then  $(G, q)$  is congruent to  $(G, p)$ .

A framework  $(G, p)$  is said to be *generic* if the set containing the coordinates of all its points is algebraically independent over the rationals. It is known that rigidity as

well as global rigidity are generic properties of  $d$ -dimensional frameworks for all  $d$ , that is, the (global) rigidity of a generic realization of a graph  $G$  depends only on the graph  $G$  and not the particular realization. We say that the graph  $G$  is *rigid*, respectively *globally rigid*, in  $\mathbb{R}^d$  if every (or equivalently, if some) generic realization of  $G$  in  $\mathbb{R}^d$  is rigid, respectively globally rigid. Many of the (global) rigidity properties of a generic framework  $(G, p)$  are determined by an associated matroid, the  *$d$ -dimensional rigidity matroid*  $\mathcal{R}_d(G)$ , defined on the edge set of  $G$ . We denote the rank of  $\mathcal{R}_d(G)$  by  $r_d(G)$ .

In what follows we shall focus on the case  $d = 2$ . In this case rigidity and the rank function of the rigidity matroid are well characterized. It is known that a graph  $G = (V, E)$  is rigid in  $\mathbb{R}^2$  if and only if  $r_2(G) = 2|V| - 3$ . It is also known that the edge set of  $G$  is independent in  $\mathcal{R}_2(G)$  if and only if each subset  $X \subseteq V$  with  $|X| \geq 2$  induces at most  $2|X| - 3$  edges [7]. Lovász and Yemini [8] characterized rigid graphs in  $\mathbb{R}^2$  by providing a formula for  $r_2(G)$ , in terms of ‘thin covers’ of  $G$ . We shall use the following refinement of their result, which uses rigid components, see [1, Section 4.4]. We define a *rigid component* of a graph  $G = (V, E)$  to be a maximal rigid subgraph of  $G$ . By the *glueing lemma* (see [10, Lemma 3.1.4]), which says that the union of two rigid graphs with at least two vertices in common is rigid, it follows that any two rigid components of  $G$  intersect in at most one vertex. Thus their vertex sets form a special ‘thin cover’ of  $G$ .

**Theorem 2.1.** [1, 8] *Let  $H = (V, E)$  be a graph with rigid components  $H_1, H_2, \dots, H_t$ . Then*

$$r_2(H) = \sum_{i=1}^t (2|C_i| - 3),$$

where  $C_i = V(H_i)$ ,  $1 \leq i \leq t$ .

Let  $G = (V, E)$  be a graph. For a family  $\mathcal{F}$  of pairwise disjoint subsets of  $V$  let  $E_G(\mathcal{F})$  denote the set, and  $e_G(\mathcal{F})$  the number, of edges of  $G$  connecting distinct members of  $\mathcal{F}$ . For a partition  $\mathcal{P}$  of  $V$  let

$$\text{def}_G(\mathcal{P}) = 3(|\mathcal{P}| - 1) - 2e_G(\mathcal{P})$$

denote the *deficiency* of  $\mathcal{P}$  in  $G$  and let

$$\text{def}(G) = \max\{\text{def}_G(\mathcal{P}) : \mathcal{P} \text{ is a partition of } V\}.$$

We say that a partition  $\mathcal{P}$  of  $V$  is *tight* if  $\text{def}_G(\mathcal{P}) = \text{def}(G)$  holds. Note that  $\text{def}(G) \geq 0$ , since  $\text{def}_G(\{V\}) = 0$ . For example, the graph  $G$  on Figure 1 has  $\text{def}(G) = 1$ . The vertex sets of the four disjoint copies of ‘ $K_4$  minus an edge’ in  $G$  form a tight partition of  $G$ .

The following rank formula (which is implicit in [6]) shows that the ‘degree of freedom’ of  $L(G)$  is equal to the deficiency of  $G$ .

**Theorem 2.2.** *Let  $G = (V, E)$  be a graph with minimum degree at least two. Then*

$$r_2(L(G)) = 2|E| - 3 - \text{def}(G). \quad (1)$$

**Proof:** First we prove that the right hand side is an upper bound on  $r_2(L(G))$ . Since  $|V(L(G))| = |E|$ , we have  $r_2(L(G)) \leq 2|E| - 3$ . Thus we may assume that  $\text{def}(G) \geq 1$ . Let  $\mathcal{Q} = \{Q_1, Q_2, \dots, Q_t\}$  be a tight partition of  $V$ . Since  $\text{def}(G) \geq 1$ , we must have  $t \geq 2$ .

For  $v \in V$  let  $B(v)$  denote the set of vertices in  $L(G)$  corresponding to the edges incident with  $v$  in  $G$ . Since  $G$  has minimum degree at least two, we have  $|B(v)| \geq 2$  for all  $v \in V$ . Let  $X_i = \cup_{v \in Q_i} B(v)$ , for  $1 \leq i \leq t$ . Since each set  $B(v)$  contains at least two vertices, we have  $|X_i| \geq 2$  for  $1 \leq i \leq t$ . Furthermore,  $|\{X_i : v_e \in X_i\}| \leq 2$  for each vertex  $v_e$  of  $L(G)$  with equality if and only if  $e \in E_G(\mathcal{Q})$ . Thus  $\sum_{i=1}^t |X_i| = |E| + e_G(\mathcal{Q})$ . Since every edge of  $L(G)$  is induced by some  $X_i$  and each set  $X \subseteq V(L(G))$  with  $|X| \geq 2$  induces at most  $2|X| - 3$  independent edges in  $\mathcal{R}_2(L(G))$ , we can deduce that

$$\begin{aligned} r_2(L(G)) &\leq \sum_{i=1}^t (2|X_i| - 3) = 2|E| + 2e_G(\mathcal{Q}) - 3t \\ &= 2|E| - 3 - \text{def}(G). \end{aligned}$$

To prove that equality holds consider the rigid components  $H_1, H_2, \dots, H_t$  of  $L(G)$  and let  $C_i = V(H_i)$  for  $1 \leq i \leq t$ . Since each set  $B(v)$ ,  $v \in V$ , induces a complete (and hence rigid) subgraph in  $L(G)$ , we must have  $B(v) \subseteq C_i$  for some  $1 \leq i \leq t$ . Furthermore, since  $|B(v)| \geq 2$  for all  $v \in V$ , the maximality of the  $C_i$ 's and the glueing lemma imply that each  $B(v)$  is contained in exactly one set  $C_i$ . Let  $Q_i = \{v \in V : B(v) \subseteq C_i\}$ ,  $1 \leq i \leq t$ . Observe that  $Q_i \neq \emptyset$  for all  $1 \leq i \leq t$ , since each rigid component  $H_i$  has at least one edge, say  $v_e v_f$ . Hence there is a vertex  $x \in V$  which is a common end-vertex of edges  $e, f$  in  $G$ . Thus  $|B(x) \cap C_i| \geq 2$  and hence, by the glueing lemma,  $B(x) \subseteq C_i$  and  $x \in Q_i$  must hold. It follows that  $\mathcal{Q} = \{Q_1, Q_2, \dots, Q_t\}$  is a partition of  $V$ .

**Claim 2.3.**  $v_e \in C_i \cap C_j$  for some  $v_e \in V(L(G))$  and  $1 \leq i < j \leq t$  if and only if  $e \in E_G(Q_i, Q_j)$ .

**Proof:** First suppose  $v_e \in C_i \cap C_j$ . Consider an edge  $v_e v_f \in E(H_i)$ . As above, we may deduce that there is a vertex  $x \in V$ , incident with  $e, f$ , with  $x \in Q_i$ . Similarly, by considering an edge  $v_e v_h \in E(H_j)$  we obtain that there is a vertex  $y \in V$ , incident with  $e, h$ , with  $y \in Q_j$ . This implies that  $e = xy$  and  $e \in E_G(Q_i, Q_j)$ .

Conversely, suppose that  $e = xy \in E_G(Q_i, Q_j)$ . Then  $B(x) \subseteq C_i$ ,  $B(y) \subseteq C_j$ . Since  $v_e \in (B(x) \cap B(y))$ , we have  $v_e \in C_i \cap C_j$ , as required. •

By using Theorem 2.1 and Claim 2.3 we obtain

$$\begin{aligned} r_2(L(G)) &= \sum_{i=1}^t (2|C_i| - 3) = 2|E| + 2e_G(\mathcal{Q}) - 3t \\ &= 2|E| - 3 - \text{def}(\mathcal{Q}) \geq 2|E| - 3 - \text{def}(G), \end{aligned}$$

which completes the proof. •

The lower bound on the minimum degree of  $G$  in Theorem 2.2 cannot be weakened. This follows by observing that if  $G$  is a star then  $L(G)$  is rigid but  $G$  is highly deficient.

### 3 Globally rigid zeolites

Globally rigid graphs in  $\mathbb{R}^2$  have been characterized by Jackson and Jordán [5]. We say that a graph  $G$  is *redundantly rigid* in  $\mathbb{R}^2$  if  $G - e$  is rigid in  $\mathbb{R}^2$  for all  $e \in E(G)$ .

**Theorem 3.1.** [5] *Let  $H$  be a graph. Then  $H$  is globally rigid in  $\mathbb{R}^2$  if and only if  $H$  is a complete graph on at most three vertices or  $H$  is 3-vertex-connected and redundantly rigid in  $\mathbb{R}^2$ .*

If  $H$  is a line graph of a 3-regular graph then a simpler characterization follows from the next theorem.

**Theorem 3.2.** *Let  $G = (V, E)$  be a 3-regular graph. Then  $L(G)$  is 3-vertex-connected and redundantly rigid in  $\mathbb{R}^2$  if and only if  $G$  is 3-edge-connected.*

**Proof:** First suppose that  $G - F$  has two connected components  $D_1, D_2$  for some  $F \subseteq E$  with  $|F| \leq 2$ . Since  $G$  is 3-regular, there must be an edge in  $D_i$  for  $i = 1, 2$ . This implies that the vertex set in  $L(G)$  corresponding to  $F$  is a separating vertex set in  $L(G)$ . Thus  $L(G)$  is not 3-vertex-connected. This proves the ‘only if’ direction.

To see the ‘if’ part, suppose that  $G$  is 3-edge-connected. This implies that  $L(G)$  is 3-vertex-connected, since each separating vertex set in  $L(G)$  gives rise to a separating edge set of  $G$  of the same size.

Next we show that  $L(G)$  is redundantly rigid. We need the following claim.

**Claim 3.3.** *Let  $H$  be a graph with minimum degree at least two and suppose that  $H$  can be made 3-edge-connected by adding at most one edge. Then  $L(H)$  is rigid.*

**Proof:** By Theorem 2.2 it suffices to show that  $\text{def}(H) = 0$ . Consider a partition  $\mathcal{P} = \{X_1, X_2, \dots, X_t\}$  of  $V(H)$  with  $t \geq 2$ . Since  $H$  can be made 3-edge-connected by adding at most one edge, all but at most two members  $X_i$  of  $\mathcal{P}$  satisfy  $e_H(X_i, V(H) - X_i) \geq 3$ , and all members satisfy  $e_H(X_i, V(H) - X_i) \geq 2$ . Hence

$$2e_H(\mathcal{P}) \geq 3t - 2 > 3(|\mathcal{P}| - 1).$$

Thus  $\text{def}(H) = 0$ . •

Now consider an edge  $p = v_e v_f$  of  $L(G)$ . This edge corresponds to a pair of edges  $e = xy, f = xz$  in  $G$  with a common end-vertex. Since  $G$  is 3-edge-connected, we can apply Claim 3.3 to  $H = G - e$  to deduce that  $L(H)$  is rigid.

It is easy to check that  $L(G) - p$  can be obtained from  $L(H)$  by adding a new vertex and connecting it to three distinct vertices of  $L(H)$ . This operation is known to preserve rigidity (in fact, connecting the new vertex to two vertices of  $L(H)$  would already preserve rigidity, see e.g. [10, Lemma 2.1.3]). Thus  $L(G) - p$  is rigid. This proves that  $L(G)$  is redundantly rigid, as required. •

By Theorems 3.1 and 3.2 we obtain:

**Corollary 3.4.** *A two-dimensional combinatorial zeolite is globally rigid if and only if its underlying graph is 3-edge-connected.*

## 4 Concluding remarks

The characterization of global rigidity provided by Corollary 3.4 has algorithmic implications, too. Given a 3-regular graph  $G$  on  $n$  vertices, the best known running time bound for testing whether  $L(G)$  satisfies both conditions of Theorem 3.1 is  $O(n^2)$ . However, testing whether  $G$  is 3-edge-connected can be done in linear time. Since  $G$  is 3-regular, this gives rise to an improved  $O(n)$  time bound for testing global rigidity of two-dimensional combinatorial zeolites.

The characterization of rigid (or globally rigid) graphs in  $\mathbb{R}^d$ , for  $d \geq 3$ , is not known. Even the special case of three-dimensional combinatorial zeolites appears to be difficult. Nevertheless, one might conjecture that for all positive integers  $d$  a  $d$ -dimensional combinatorial zeolite is globally rigid if and only if its underlying  $(d+1)$ -regular graph is  $(d+1)$ -edge-connected (or possibly  $(d+1)$ -vertex-connected). This natural extension of Corollary 3.4, which is correct for  $d \leq 2$ , fails in  $\mathbb{R}^3$ : a counterexample, due to Bill Jackson [4], is shown in Figure 2.

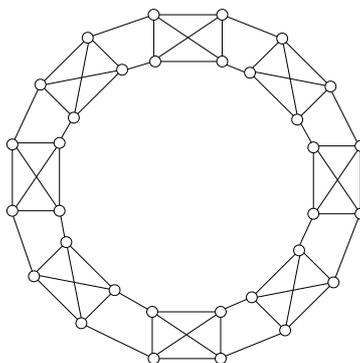


Figure 2: A 4-vertex-connected 4-regular graph  $G$  for which  $L(G)$  is not rigid (and hence not globally rigid) in  $\mathbb{R}^3$ . The line graph of  $G$  behaves like a body-and-hinge framework whose underlying graph is a cycle of length 8, and hence it is easily shown to be flexible. (Thus a cycle of length 7 would also be a counterexample.)

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